

Harshit Gangal Software Engineer

Atomic Distributed Transactions In Vitess



Manan Gupta Software Engineer



Achieving Atomic Commits across shards



- Multi Commit
- Strict Single Mode
- Consistent Lookup Vindex
- Need for Atomic Cross Shard Commit

- Best Effort Commit
- Commit in a sequential shard order
- Failures reported with shard information
- Application-Level Rollback/Rollforward
- Single shard transactions are ACID



- Multi Commit
- Strict Single Mode
- Consistent Lookup Vindex
- Need for Atomic Cross Shard Commit

- Restrict Transactions to Single Shard
- Rollbacks on Cross Shard



- Multi Commit
- Strict Single Mode
- Consistent Lookup Vindex
- Need for Atomic Cross Shard Commit

- Global Secondary Index
- Open Cross-shard Transaction
- Uses Locking and Transaction sequence

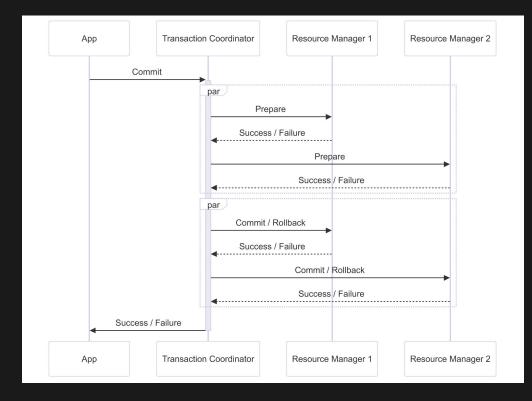
- Multi Commit
- Strict Single Mode
- Consistent Lookup Vindex
- Need for Atomic Cross Shard Commit

- Reduce Application complexity
- Handle failure modes

Two-Phase Commit

Overview

- Transaction Coordinator
- Resource Manager



Two-Phase Commit

Guarantees

- Prepare Protocol
 - Never abort a transaction, unless requested
 - Never refuse a commit
 - After a crash, reinstate transaction to its prepared state on recovery



Two-Phase Commit

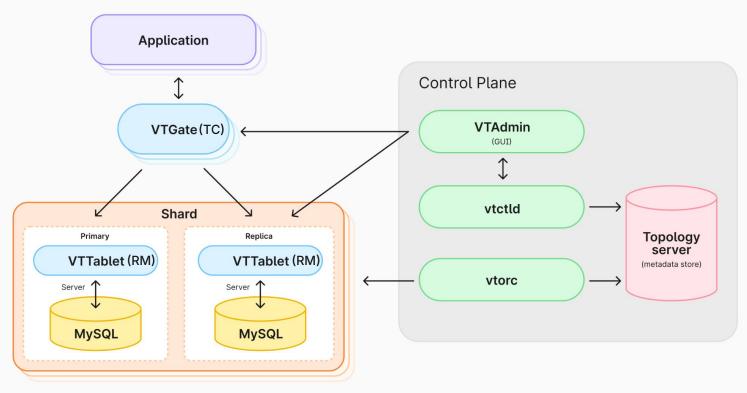
Guarantees

- Prepare Protocol
 - Never abort a transaction, unless requested
 - Never refuse a commit
 - After a crash, reinstate transaction to its prepared state on recovery
- MySQL
 - Transaction does not abort
 - High connection wait timeout
 - No TCP connection
 - Commit almost never fails
 - No group replication
 - Need for Transaction Logs for crash recovery



Vitess Architecture

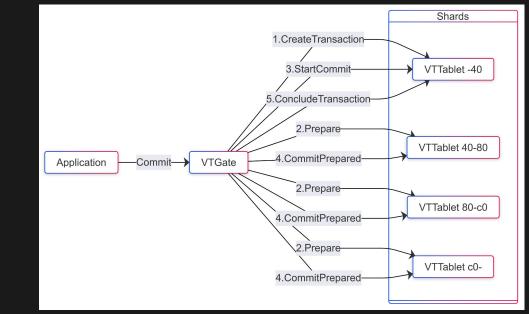
Vitess Runtime





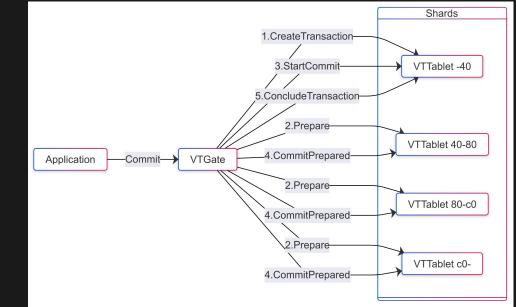
Create Transaction Record

- Generate Unique Distributed Transaction ID (DTID)
- One of Participating VTTablet takes role of Metadata Manager
- Persist the transaction metadata in a separate autocommit transaction



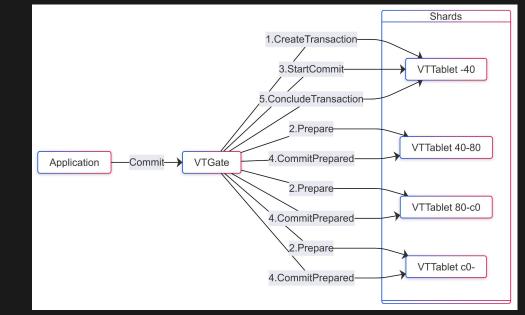
Prepare

- VTTablet persists the transaction logs in separate autocommit transaction
- Moves open transaction out of transaction timeout scope



Start Commit

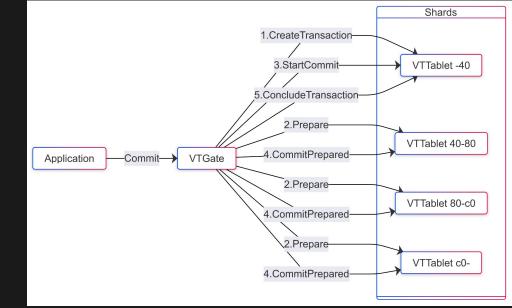
- Metadata Manager marks the transaction status to Commit on same transaction
- Commits the open transaction





Commit Prepared

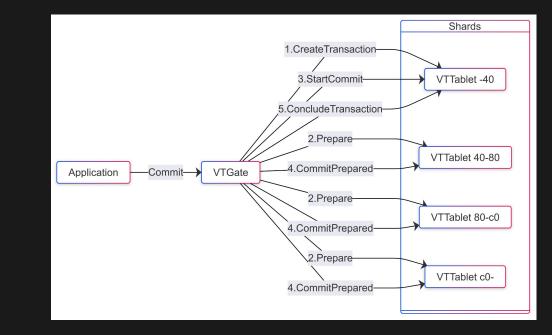
• Commits all the prepared transactions





Conclude Transaction

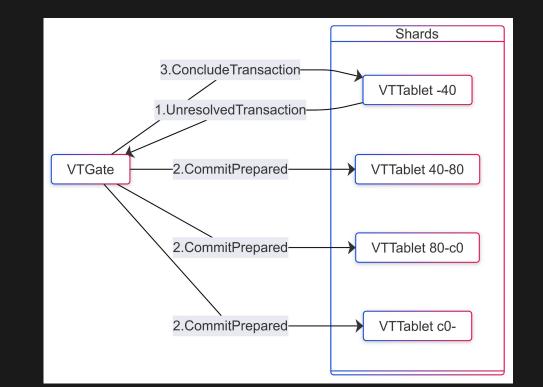
• Removes the Transaction Record





Transaction Resolution Watcher

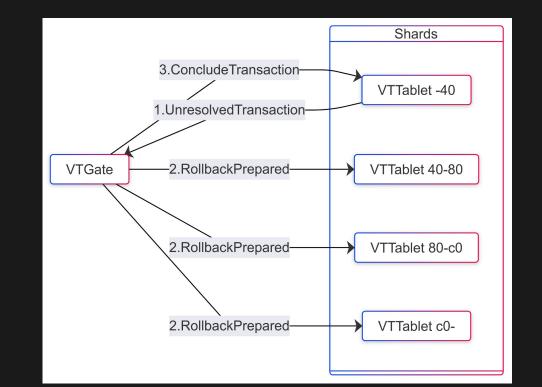
- Handling unresolved transactions
- Transaction State
 - Commit
 - Rollback
 - Prepare





Transaction Resolution Watcher

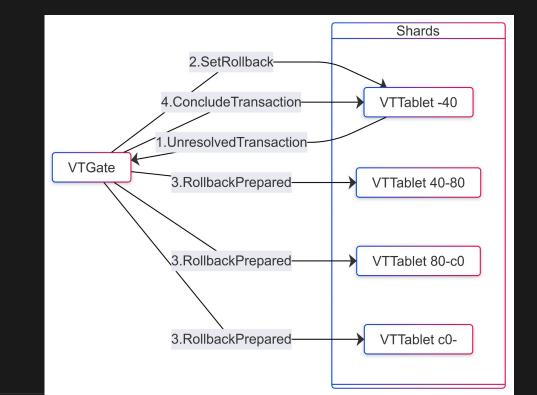
- Handling unresolved transactions
- Transaction State
 - Commit
 - Rollback
 - Prepare





Transaction Resolution Watcher

- Handling unresolved transactions
- Transaction State
 - Commit
 - Rollback
 - Prepare





Design Benefits

- Zero impact single shard transactions
- Stateless coordinator model
- Avoiding prepare phase of Metadata Manager
- Relaxed Isolation for scalability



Resilience During Disruptions





MySQL Restarts

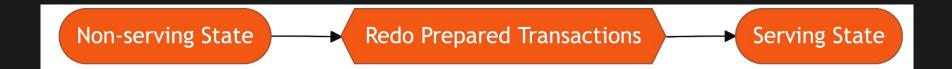
- All Ongoing Connections Dropped!
 - This includes all Prepared Transactions
- MySQL starts in Super_Read_Only state





VTTablet Restarts

- Similar to MySQL restart.
- All connections dropped
- VTTablet starts in non-serving state





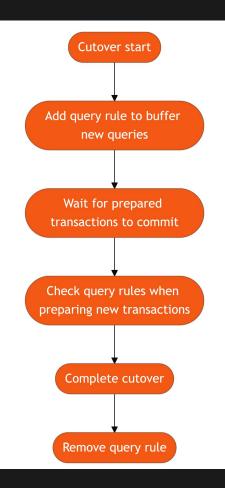
VTGate Restarts

- Nothing to do!
- Transaction resolution watcher takes care of it!



OnlineDDL

- Cannot have any prepared transactions dependent on the old schema.
- Use query rules to tie everything together.



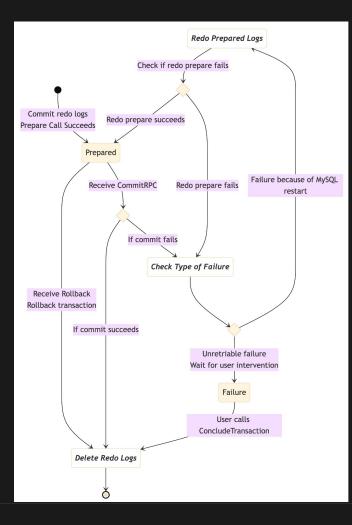


MoveTables and Reshard

- Very similar solution to OnlineDDL
- Add query rules to fail new queries and prepared transactions
- Wait for open prepared transactions
- Remove the query rule.



VTTablet State Diagram





Monitoring





Monitoring - VTAdmin UI

C G K S T

In Flight Distributed Transactions

Classic		Keyspace	Abando	on Age 🛈				
Clusters	1	customer (local) ~	5sec		~			
Gates	1							
		ID		State	Participants	Time Created		Actions
Keyspaces	2	customer:80-:1736500903351809003		COMMIT	customer/-80customer/80-	2025-01-10 3:04 PM -	+05:30	0
Schemas	5					22 seconds ago		
Tablets	9	Showing 1 of 1					Conclude Transact	ion



Monitoring - VTAdmin UI

V

← All Unresolved Transactions /

Cluster: local

customer:80-:1736500903351809003



ID			St	ate	Participants	Time Created		Action
custom	er:80-:17365	00903351809	0003 CC	DMMIT	customer/-80customer/80-	2025-01-10 3:04 PM + 42 seconds ago	05:30	()
Showing 1	of 1						Conclude Transac	tion
Shard	State	Message	Time Created	Sta	tements			
-80	PREPARED	÷	2025-01-10 3:04 PM +05:30 42 seconds ago) ins	ert into corder(order_id, custome	r_id, sku, price) values	(1010, 101, 'a',	101)



Migrations

Monitoring - MySQL Query

• In case of a commit failure, `SHOW WARNINGS` can be used to retrieve the dtid.



Monitoring - Metrics & Stats (VTGate)

- CommitUnresolved Count of transactions that failed during Commit.
- CommitModeTimings Timings for how long the commit phase takes to complete.

```
"CommitModeTimings": {"TotalCount":3,"TotalTime":32325084,"Histograms":{"Multi":
{"500000":0,"10000000":0,"5000000":0,"10000000":0,"50000000":0,"5000000"
0":0,"1000000000":0,"50000000":0,"100000000":0,"inf":0,"Count":1,"Time":10804250}
,"TwoPC":
{"500000":0,"10000000":0,"50000000":0,"10000000":1,"50000000":1,"100000000":0,"50000000
0":0,"100000000":0,"50000000":0,"10000000":1,"50000000":1,"100000000":0,"50000000
0":0,"100000000":0,"50000000":0,"10000000":1,"50000000":1,"100000000":0,"50000000
0":0,"100000000":0,"50000000":0,"10000000":1,"50000000":1,"100000000":0,"50000000
0":0,"100000000":0,"5000000":0,"10000000":1,"50000000":1,"100000000":0,"5000000
0":0,"1000000000":0,"50000000":0,"100000000":0,"inf":0,"Count":2,"Time":21520834}
}},
"CommitUnresolved": 0,
```



Monitoring - Metrics & Stats (VTTablet)

- QueryTimings Timings of individual phases of 2PC CreateTransaction, Prepare, StartCommit, SetRollback, CommitPrepared, RollbackPrepared and Resolve.
- Unresolved Gauge of number of transactions running longer than the abandon age.



Testing

Consession ((consession)) (Ball ((consession)) ther loaha [] nottekesse () (() = () =) () =) ALLER THAT BROKE Orbitan **CITILE** nobtel Leaners thosses bock hs) ((~sooik) the: _dock)~) Themechae? adhasoodha`)a(dadasthtt.tt rochcith asoonn))) bkasssfaddmcleekee deckiebte ((([_besnctchiessdeditons])) 66 obhrc. e Scischdobsassosa))()) Beater 11 spemba _stsheurshsa)(ssienseitlerer a 100 m l at (cens(L phoorn **all)sce**cTc32)ncoent)) **ense** TREESECT. DCILG ((dcx(odok)))_c=neccc(e) m. . . eklerssf (==sdoohackcinx)))) IDIAN HARDERATS (F) Lboro (ask tochel=("coop hosskoettna ((BeekEih))). setthe);shoe. V VENSOE THE Interview interview netonetense ((énentre)DD7) ttat croonweathepedasnis an ebect (: (percere))) haco CO PERMISES



Testing - Challenges

- Ensure everything works as expected.
- Disruption handling code should preserve all the 2PC guarantees
- To test the code, we need a reliable way to inject errors!



Go Build Tags To The Rescue

// <u>go:build</u> debug2PC
package vtgate
import (
"context"
"vitess.io/vitess/go/vt/callerid"
"vitess.io/vitess/go/vt/log" querypb "vitess.io/vitess/go/vt/proto/query"
vtrpcpb "vitess.io/vitess/go/vt/proto/vtrpc"
"vitess.io/vitess/go/vt/vterrors"
) 🕊

const DebugTwoPc = true no usages 🗳 Manan Gupta

// <u>go:build</u> !debug2PC
package vtgate
import (
"context"
<pre>querypb "vitess.io/vitess/go/vt/proto/query"</pre>
<pre>// This file defines debug constants that are always false. // This file is used for building production code.</pre>
<pre>// We use go build directives to include a file that defines the constant to true // when certain tags are provided while building binaries.</pre>
// This allows to have debugging code written in normal code flow without affecting
// production performance.
const DebugTwoPc = false 6 usages A Manan Gupta



Go Build Tags To The Rescue

```
txPhase = Commit2pcCreateTransaction
if err = txc.tabletGateway.CreateTransaction(ctx, mmShard.Target, dtid, participants); err \neq nil { return err }
if DebugTwoPc { // Test code to simulate a failure after RM prepare
    if terr := checkTestFailure(ctx, "TRCreated_FailNow", nil); terr ≠ nil {
        return terr
    }
}
txPhase = Commit2pcPrepare
prepareAction := func(ctx context.Context, s *vtgatepb.Session_ShardSession, logging *econtext.ExecuteLogger) error {
    if DebugTwoPc { // Test code to simulate a failure during RM prepare
        if terr := checkTestFailure(ctx, "RMPrepare_-40_FailNow", s.Target); terr \neq nil { return terr }
    7
    return txc.tabletGateway.Prepare(ctx, s.Target, s.TransactionId, dtid)
}
if err = txc.runSessions(ctx, rmShards, session.GetLogger(), prepareAction); err \neq nil { return err }
if DebugTwoPc { // Test code to simulate a failure after RM prepare
    if terr := checkTestFailure(ctx, "RMPrepared_FailNow", nil); terr ≠ nil { return terr }
}
```



Comprehensive Testing

- Unit Testing
- Disruptions handling tests using gobuild trick to inject delays and ensure guarantees are still met.
- Other miscellaneous end to end tests for metrics, UI, etc.
- Fuzzer testing for extended reliability and to find unknown cases!
 - This worked very well for us in the past with foreign keys work.



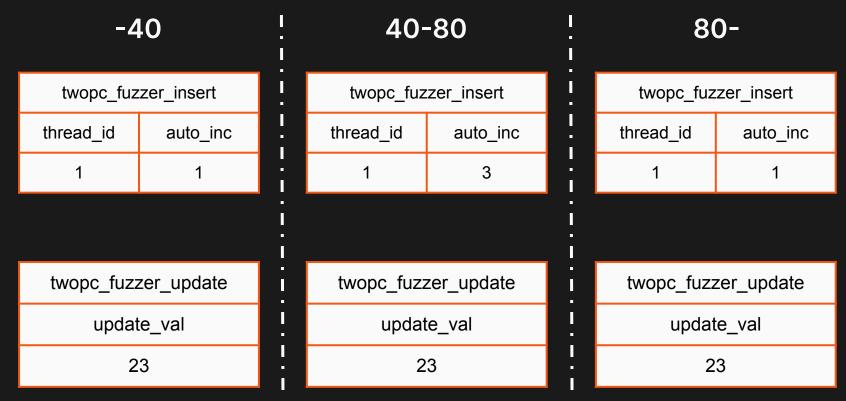
Fuzzer Testing Ideas Explained

- Multiple threads running distributed transactions.
- We want to check they're all atomic.
- Another thread running disruptions.
- twopc_fuzzer_insert and twopc_fuzzer_update tables we use.
- Each transaction inserts a row in the first table, and updates a row in the latter table in all the shards.
- For the insertion, we have an auto-increment column. For the update, we have a column that we increment with a random value, but it is same across all the shards.
- Check atomiticity by ensuring that column in twopc_fuzzer_update matches in all shards.
- Check order of commit by ensuring that for each thread, the order of auto increment column matches in twopc_fuzzer_insert in all the shards.











40-80 -40 80twopc_fuzzer_insert twopc_fuzzer_insert twopc_fuzzer_insert Π thread id auto_inc thread id auto_inc thread id auto_inc 3 1 1 1 1 2 2 2 3 2 4 twopc_fuzzer_update twopc_fuzzer_update twopc_fuzzer_update update_val update_val update_val Π 123 123 123



40-80 -40 80twopc_fuzzer_insert twopc_fuzzer_insert twopc_fuzzer_insert Π thread id auto_inc thread id auto_inc thread id auto_inc 3 1 1 1 1 2 2 2 4 twopc_fuzzer_update twopc_fuzzer_update twopc_fuzzer_update update_val update_val update_val Π 123 123 23



40-80 -40 80twopc_fuzzer_insert twopc_fuzzer_insert twopc_fuzzer_insert Π thread id auto_inc thread id auto_inc thread id auto_inc 2 3 1 1 1 2 2 2 3 4 1 twopc_fuzzer_update twopc_fuzzer_update twopc_fuzzer_update update_val update_val update_val Π 123 123 123



Future Enhancements

- Read Isolation Guarantee
- Distributed Deadlock Avoidance



